

Hall-D Software

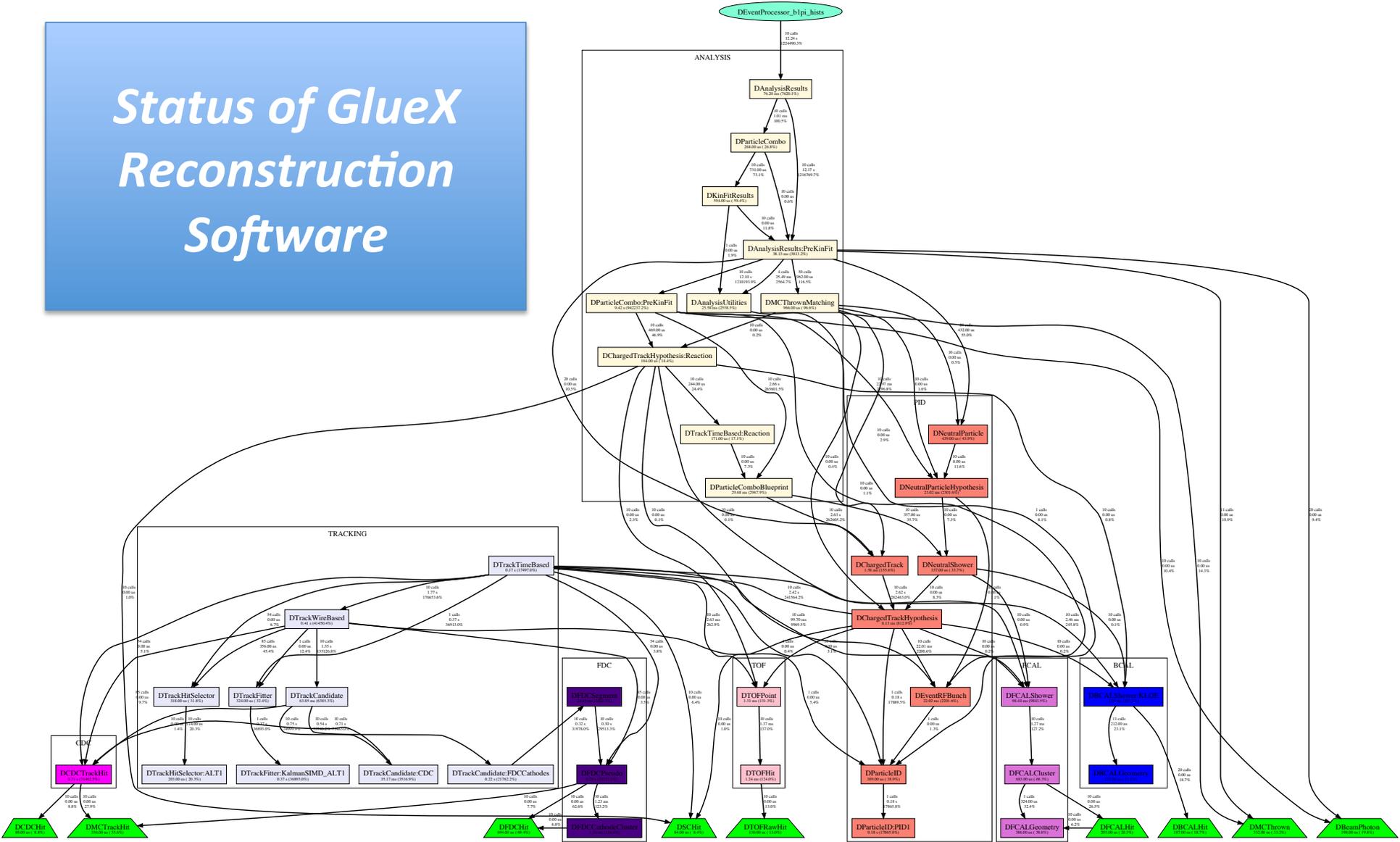
David Lawrence

June 7, 2012

What is JANA?

- JLab **ANA**lysis framework (more accurately, a *reconstruction*) framework
- C++ framework that formalizes the organization of algorithms and data transfer for event based processing
- Multi-threaded event processing
- Numerous additional features:
 - Configuration parameters
 - Web-based Resource retrieval
 - Plugins
 - Automatic ROOT tree creation
 - Calibration DB API

Status of GlueX Reconstruction Software



View Controls

ZOOM

Transverse Coordinates
 x/y
 r/phi

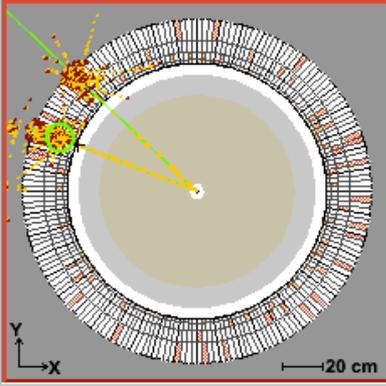
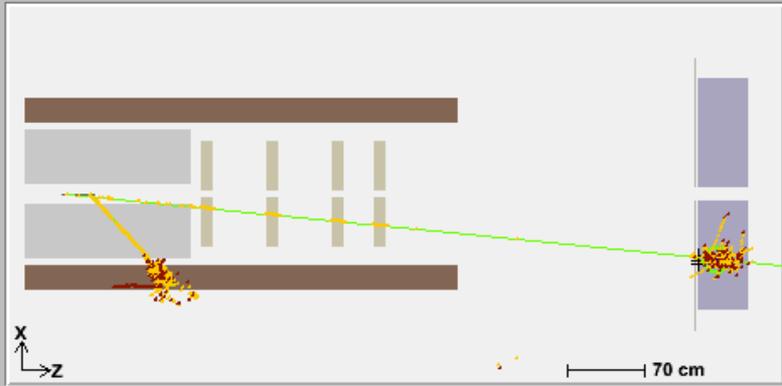
Event Controls

continuous
 delay:

Info

Run: -----
 Event: 5

Inspectors

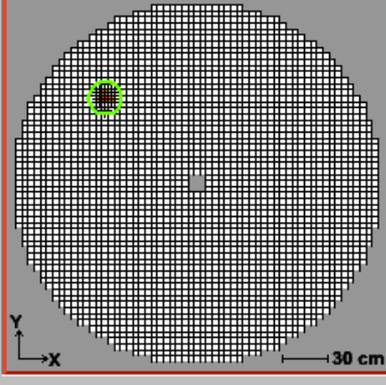
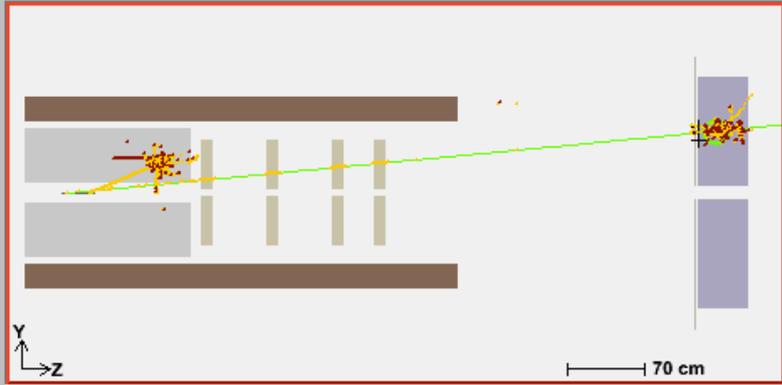


BCAL colors

- 20,0 MeV
- 15,8 MeV
- 12,2 MeV
- 9,0 MeV
- 6,3 MeV
- 4,0 MeV
- 2,3 MeV
- 1,1 MeV
- 0,3 MeV

Track Draw Options

DTrackCandidate: <default>
 DTrackWireBased: <default>
 DTrackTimeBased: <default>
 DChargedTrack: <default>
 DNeutralParticle
 DMCThrown
 DMCTrajectoryPoint



Debugger

FCAL colors

- 200,0 MeV
- 111,7 MeV
- 62,3 MeV
- 34,8 MeV
- 19,4 MeV
- 10,9 MeV
- 6,1 MeV
- 3,4 MeV
- 1,9 MeV

Hit Draw Options

CDC
 CDC Drift Time
 CDCTruth
 FDC Wire
 FDC Pseudo
 FDCTruth
 TOF
 TOFTruth
 FCAL
 BCAL

Track Info

Thrown

trk:	type:	p:	theta:	phi:	z:
1	gamma	1	8	2,376	54,88
2	gamma	1	50	2,768	75,63
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2/11/14

Reconstructed

trk:	type:	p:	theta:	phi:	z:	chisq/Ndof:	Ndof:	FOM:	cand:
1	gamma	0,9706	45,43	2,766	65	N/A	N/A	N/A	-----
2	gamma	0,8723	8,161	2,387	65	N/A	N/A	N/A	-----
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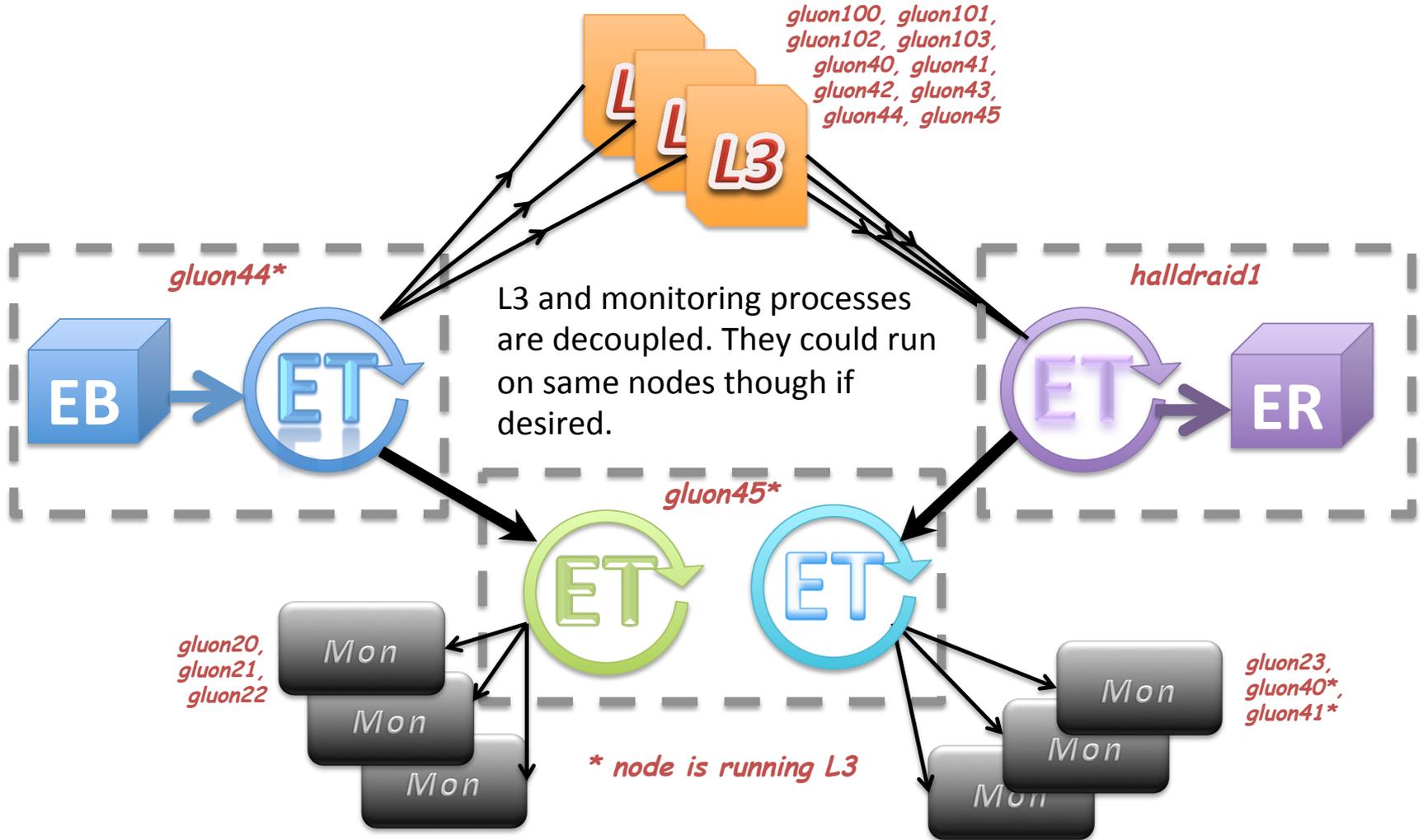
cand:

Distributed Computing in GlueX

- Online systems
 - Monitoring farm (*ET*)
 - L3 trigger farm (*ET*)
- Offline systems
 - Raw data reconstruction/analysis (*Auger/PBS*)
 - Simulation (*Open Science Grid/Auger/PBS*)

L3 and monitoring architecture

for 2013 Online Data Challenge



*n.b. all L3 machines connected via InfiniBand

Jefferson Lab Workflow Software System

Feature Requirements

Jie Chen, John Goetz, Vardan Gyurjyan, Mark Ito, Chris Larrieu,
David Lawrence, Sandy Philpott, Chip Watson, and Dennis Weygand

November 6, 2013

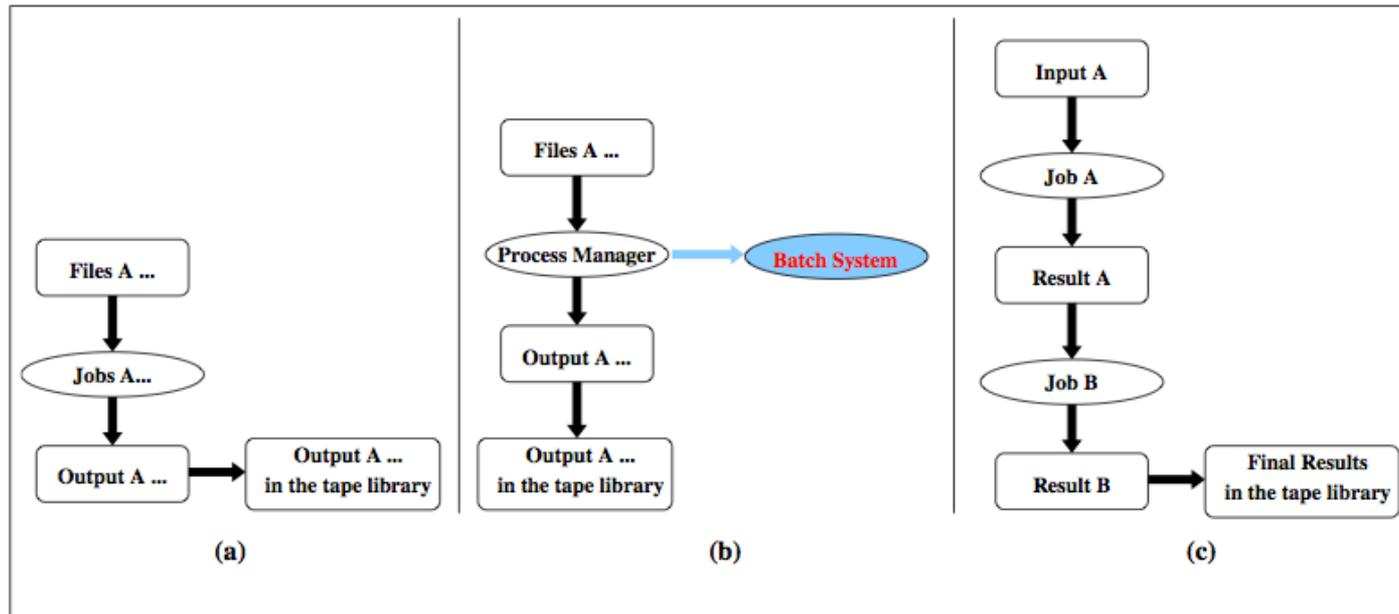
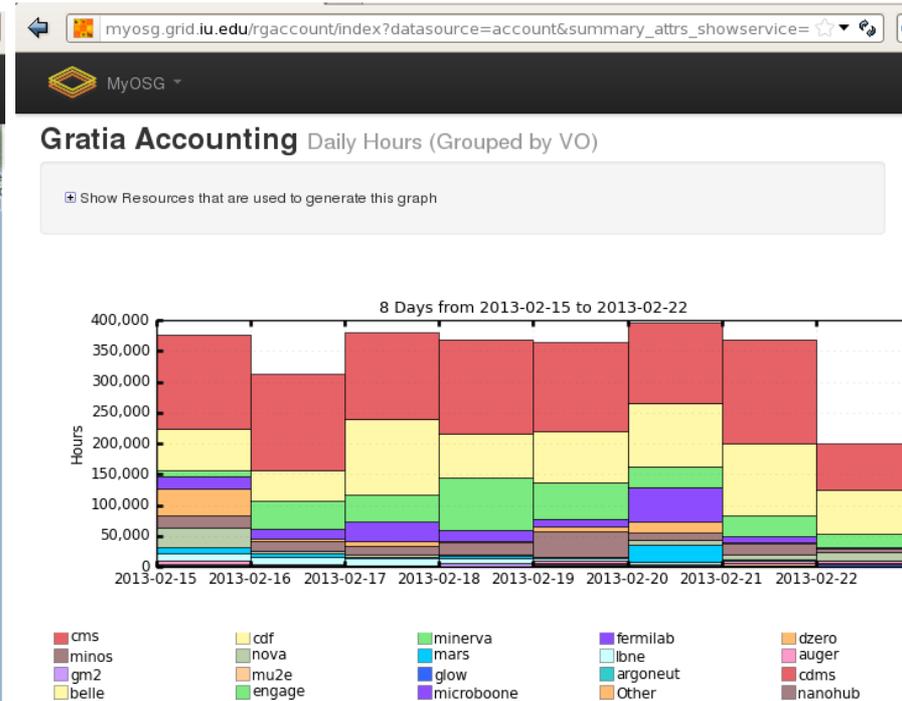
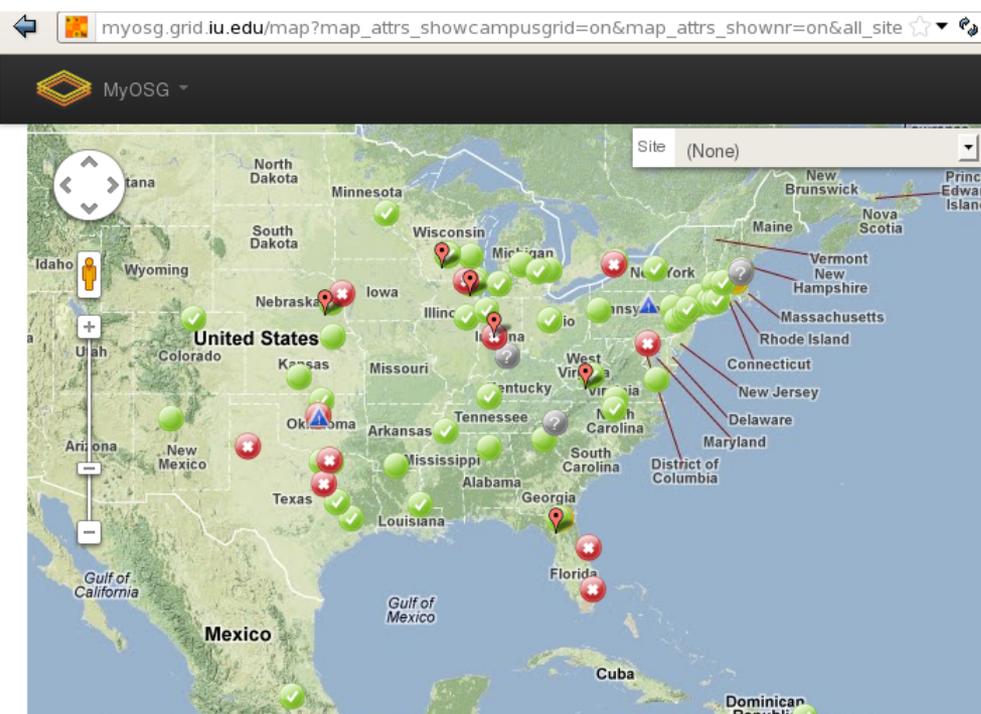


Figure 1: Three different types of workflow.

In the figure, type (a) presents a simplest workflow that a set of N jobs process a set of M data files where N and M can be 1. There is no dependency among the jobs. The output of the jobs appear on disk and some or all of the output files eventually end up in the tape library. Type (b) presents a similar workflow that is partially managed by an external process management system such as CLARA. Finally, type (c) workflow happens less often and it expresses a multi-staged data analysis, which consists a series of jobs each of which cannot start until the previous job finishes or the output files from the previous job are available.

OSG Context

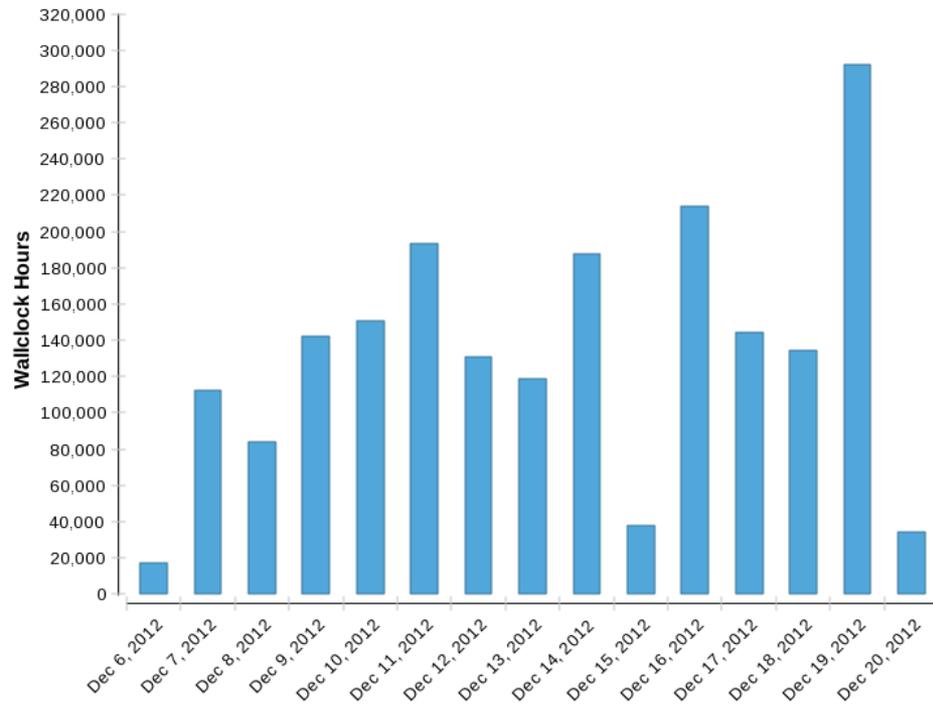
- Open Science Grid – founded 2004
- primary driver supporting LHC experiments in N/S America
- over 75,000 cores, running a distribution of Linux
- sites at 72 institutions including 42 universities, 90 sites (US, Brazil)
- centrally managed and operated by full-time staff (GOC @ I.U.)



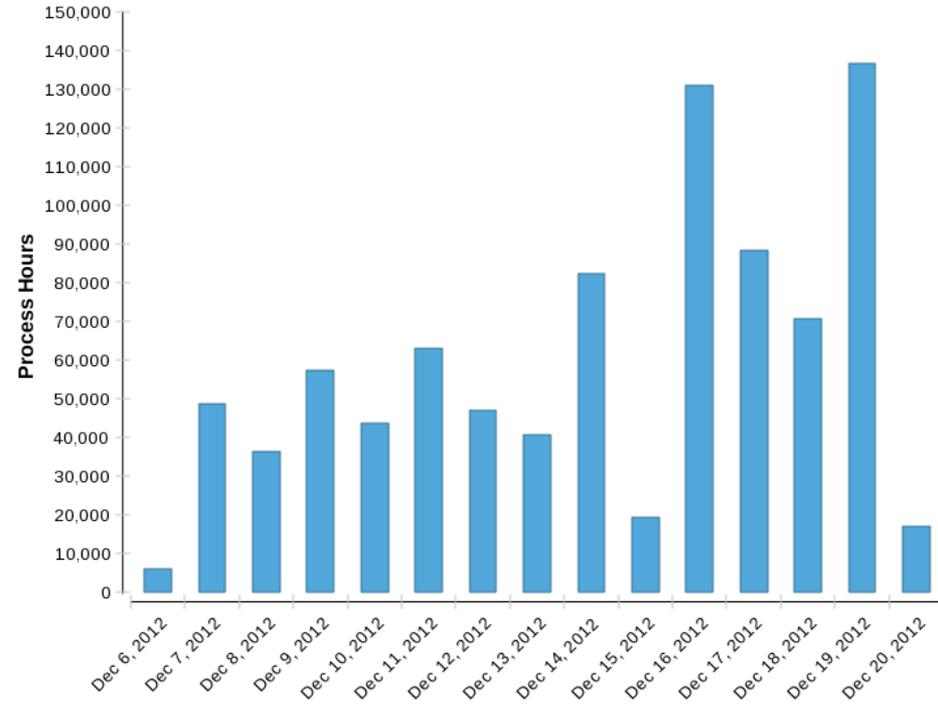
Achievements

- cpu availability was very high (>10,000 cores peak)
- production efficiency was not great (40 – 60%)
- part of inefficiency is due to pre-emption (opportunistic)
- understanding sources of inefficiency is reason why we stopped @5B events

Daily Usage by VO (Wallclock Hours)

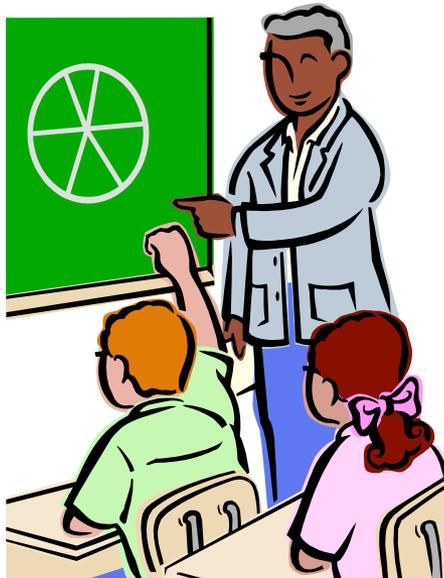


Daily Usage by VO (Process Hours)

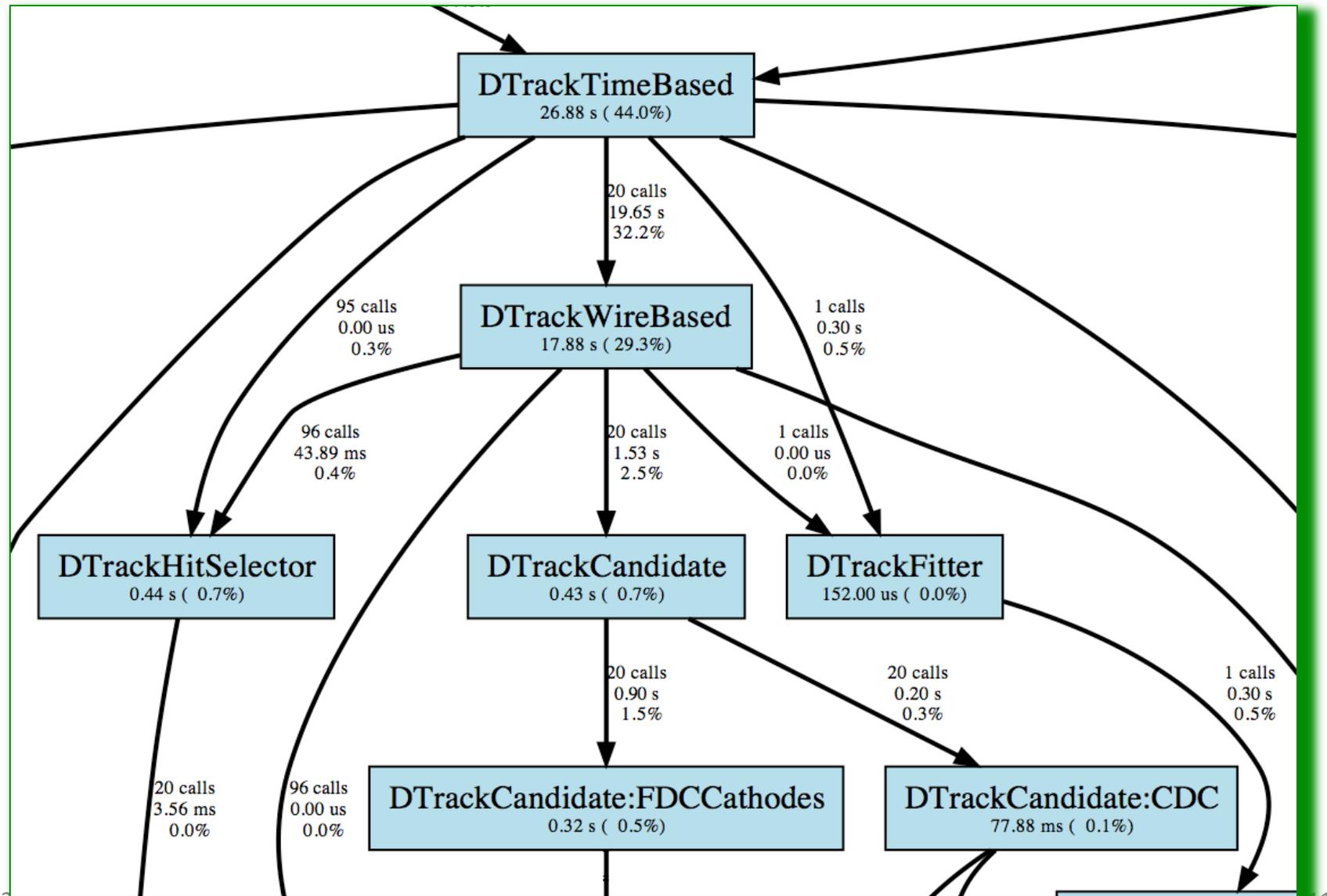


Richard Jones
GlueX Collaboration Meeting,
Newport News, Feb. 21-23,
2013

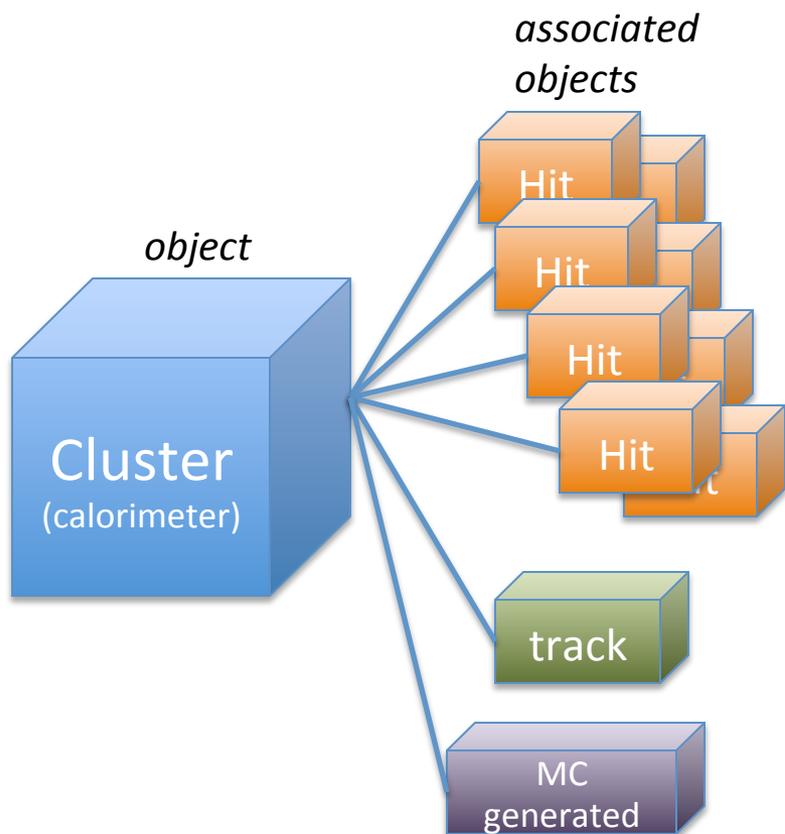
Backup Slides



A closer look at *janadot*



Associated Objects



- A data object may be associated with any number of other data objects having a mixture of types
- Each data object has a list of “associated objects” that can be probed using a similar access mechanism as for event-level object requests

```
vector<const DCluster*> clusters;  
loop->Get(clusters);  
for(uint i=0; i<clusters.size(); i++)  
{  
    vector<const DHit*> hits;  
    clusters[i]->Get(hits);  
    // Do something with hits ...  
}
```

Configuration Parameters

in a factory's init method one might write ...

```
MIN_SEED_HITS = 4;  
MAX_STEP_SIZE = 3.0; // cm
```

*Variables are data
members of factory class*

*Value may be overwritten if user
specifies a value at run time*

```
gPARMS->SetDefaultParameter("TRKFIND:MIN_SEED_HITS", MIN_SEED_HITS);  
gPARMS->SetDefaultParameter("TRK:MAX_STEP_SIZE" , MAX_STEP_SIZE  
    , "Maximum step size in cm to take when swimming a track with adaptive step sizes");
```

***NEW:** Optional 3rd argument allows
description to be stored with parameter*

- Parameters can be set via command line or configuration file
 - Complete list of parameters can be dumped using option `--dumpconfig`
 - Parameters can read in using option `--config=filename`

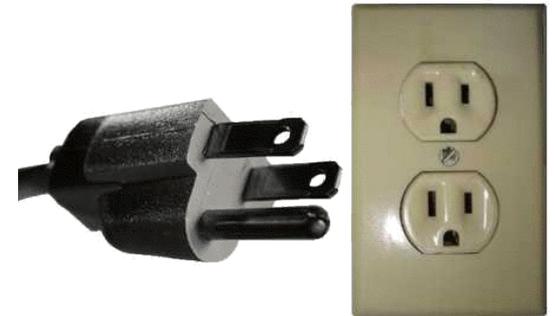
Example Configuration Parameter dump file

```
#
# JANA Configuration parameters (auto-generated)
#
# created: Wed May  5 11:32:54 2010
# command: hd_ana --dumpconfig -PEVENTS_TO_KEEP=1 --auto_activate=DChargedTrack hdgeant_smeared.hddm
#

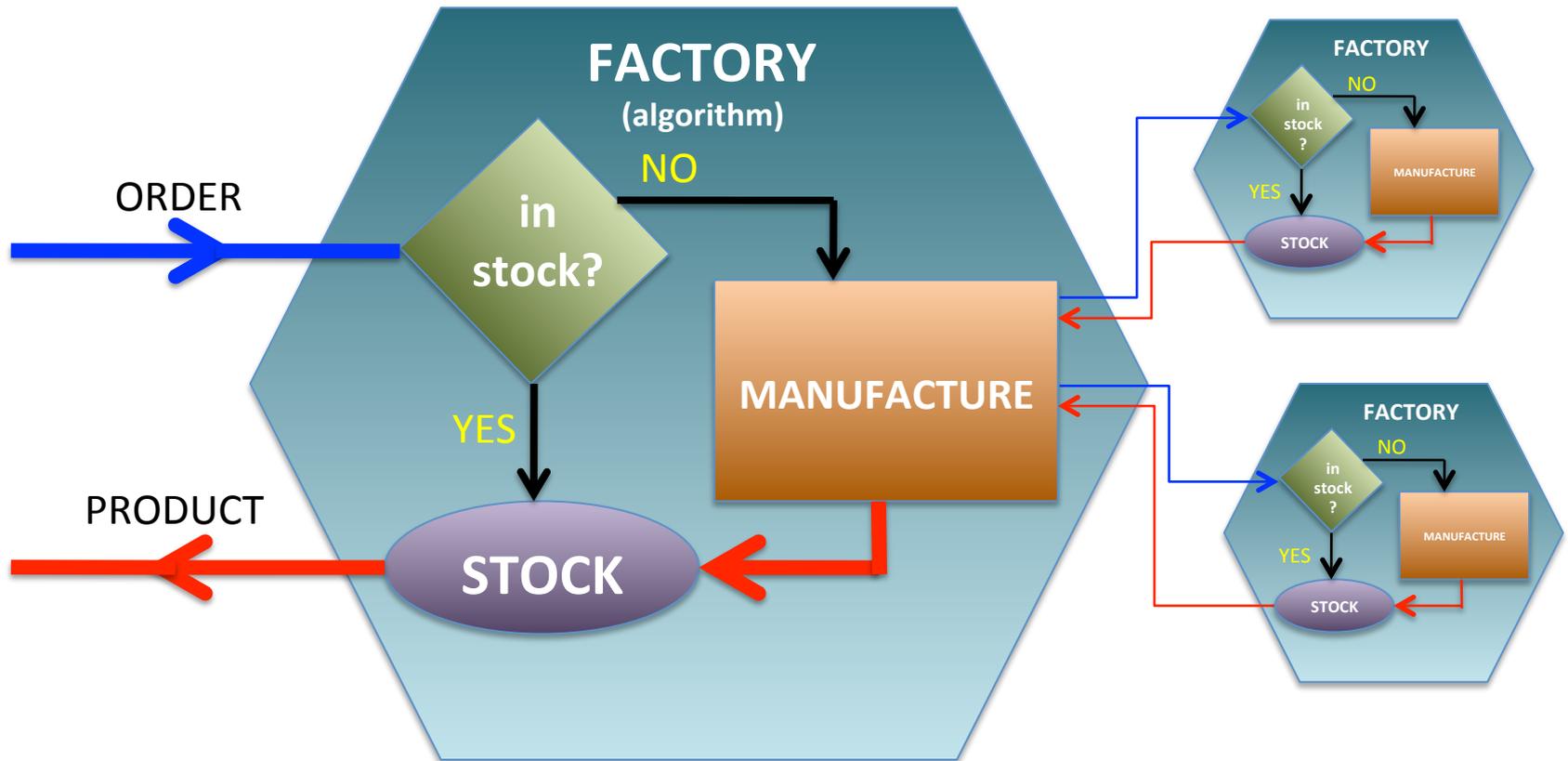
BCALRECON:BREAK_THRESH_TRMS          5
BCALRECON:CLUST_THRESH                0.02
BCALRECON:MERGE_THRESH_DIST          40
BCALRECON:MERGE_THRESH_TIME          2.5
BCALRECON:MERGE_THRESH_XYDIST        40
BCALRECON:MERGE_THRESH_ZDIST         30
BCALRESPONSE:CELL_THRESHOLD_OUTER    0.001
BCALRESPONSE:CROSS_TALK_PROB         0.03
BCALRESPONSE:DARK_RATE_GHZ           0.041
BCALRESPONSE:DEVICE_PDE               0.12
BCALRESPONSE:FADC_WINDOW_NS           100
BCALRESPONSE:OCCUPANCY_FRACTION_LIMIT 0.05
BCALRESPONSE:PHOTONS_PER_SIDE_PER_MEV_IN_FIBER 75
BCALRESPONSE:SAMPLING_COEF_A          0.042
BCALRESPONSE:SAMPLING_COEF_B          0.013
BCALRESPONSE:SAMPLING_FRACTION        0.15
BCALRESPONSE:TIMESMEAR_COEF_A         0.0989949
BCALRESPONSE:TIMESMEAR_COEF_B         0
BFIELD_MAP                            Magnets/Solenoid/solenoid_1500_poisson_20090814_01
BFIELD_TYPE                            CalibDB
CDC:Z_MAX                              167
CDC:Z_MIN                              17
EVENTS_TO_KEEP                          1 # Maximum number of events for which event processors are cal
EVENTS_TO_SKIP                          0 # Number of events that will be read in WITHOUT calling event
FCAL:BUFFER_RADIUS                      8
FCAL:FCAL_CRITICAL_ENERGY               0.035
FCAL:FCAL_RADIATION_LENGTH              3.1
FCAL:FCAL_SHOWER_OFFSET                 1
FCAL:MIN_CLUSTER_BLOCK_COUNT            2
FCAL:MIN_CLUSTER_SEED_ENERGY            0.035
FCAL:NON_LIN_COEF_A1                    0.53109
FCAL:NON_LIN_COEF_A2                    0.463044
FCAL:NON_LIN_COEF_alfa1                 1.01919
FCAL:NON_LIN_COEF_alfa2                 1.03614
FCAL:NON_LIN_COEF_B1                    2.66426
FCAL:NON_LIN_COEF_B2                    2.4628
FCAL:NON_LIN_COEF_C1                    2.70763
FCAL:NON_LIN_COEF_C2                    2.39377
FCAL:RHG_RADIUS                         30
GEOM:ENABLE_BOUNDARY_CHECK              1 # Enable boundary checking (superceeds any setting in DReferer
GEOM:MAX_BOUNDARY_SEARCH_STEPS           30 # Maximum number of steps (cells) to iterate when searching f
JANA:JERR_TAG                            JANA ERROR>>
JANA:JERR_THREADSTAMP_FLAG              0 # string prefixed to all lines sent to jerr ofstream
# if non zero prepend nthread id to each message printed to
```

Plugins

- JANA supports plugins: pieces of code that can be attached to existing executables to extend or modify its behavior
- Plugins can be used to add:
 - Event Processors
 - Event sources
 - Factories (additional or replacements)
- Examples:
 - Plugins for creating DST skim files
 - Reconstruction is done once with output to multiple files
 - `hd_ana --PPLUGINS=kaon_skim,ppi+pi-_skim run012345.evio`
 - Plugins for producing subsystem histograms
 - Single ROOT file has histograms from several pieces of code
 - `hd_root --PPLUGINS=bcal_hists,cdc_hists,tof_hists ET:GlueX`

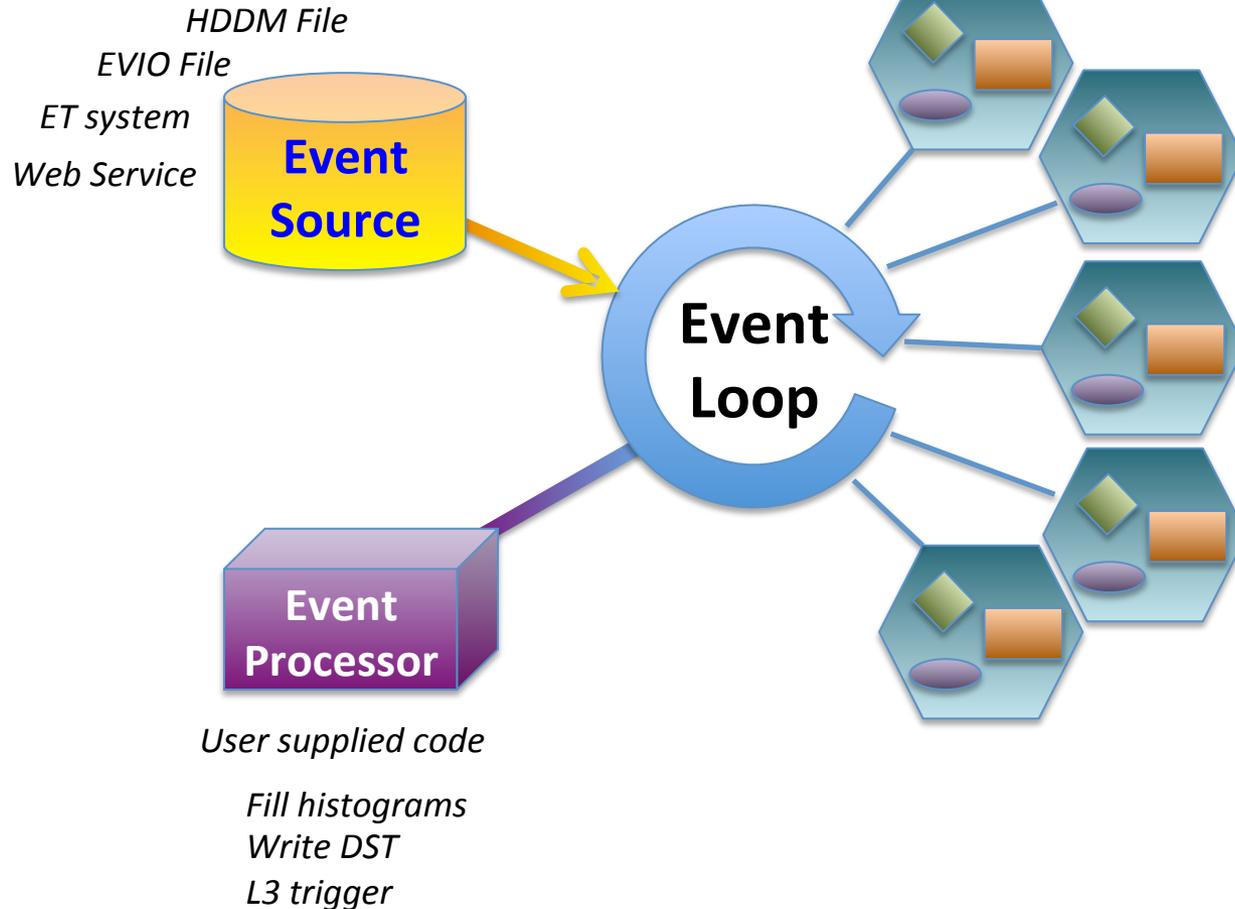


Factory Model



Data on demand = Don't do it unless you need it
Stock = Don't do it twice — Conservation of CPU cycles!

Complete Event Reconstruction



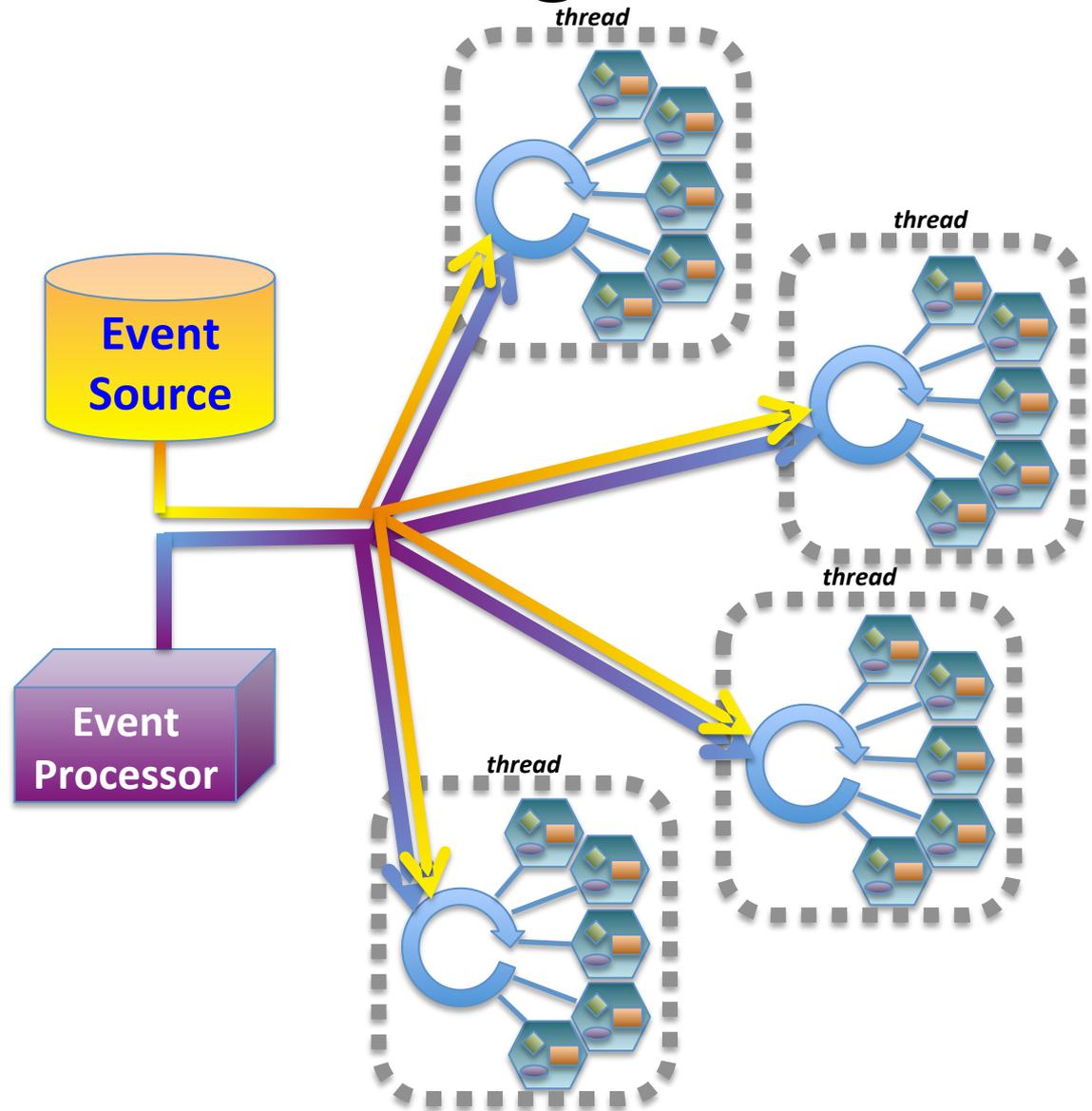
Framework has a layer that directs object requests to the factory that completes it

Multiple algorithms (factories) may exist in the same program that produce the same type of data objects

This allows the framework to easily redirect requests to alternate algorithms specified by the user at run time

Multi-threading

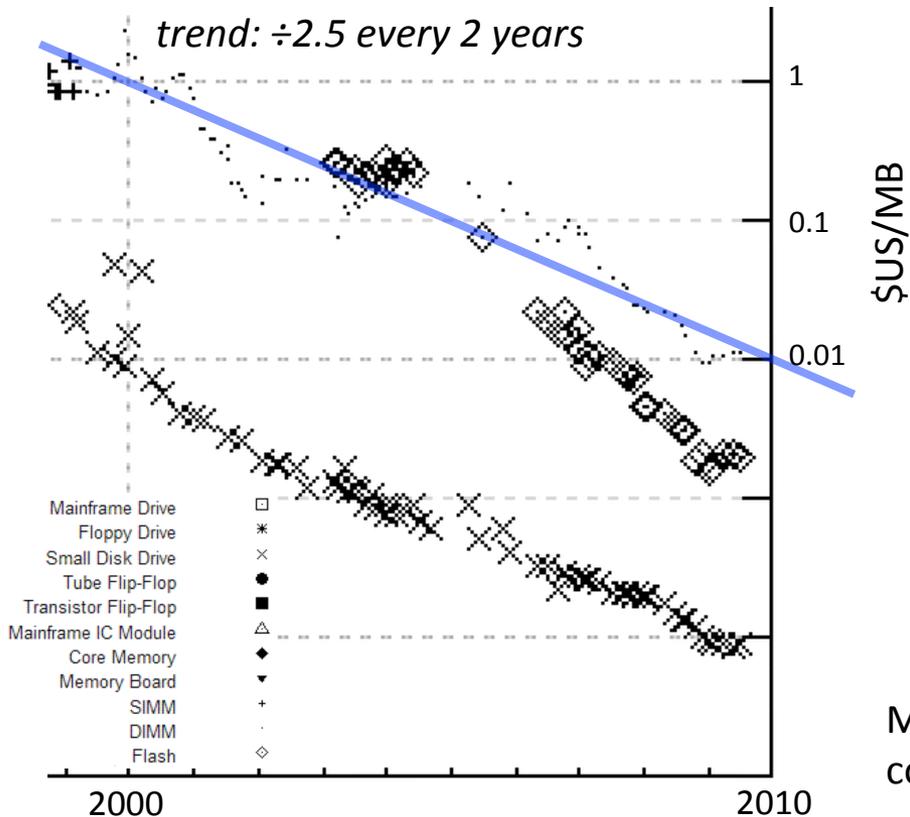
- *Each thread has a complete set of factories making it capable of completely reconstructing a single event*
- *Factories only work with other factories in the same thread eliminating the need for expensive mutex locking within the factories*
- *All events are seen by all Event Processors (multiple processors can exist in a program)*



Memory prices for last 2 decades

Memory Prices

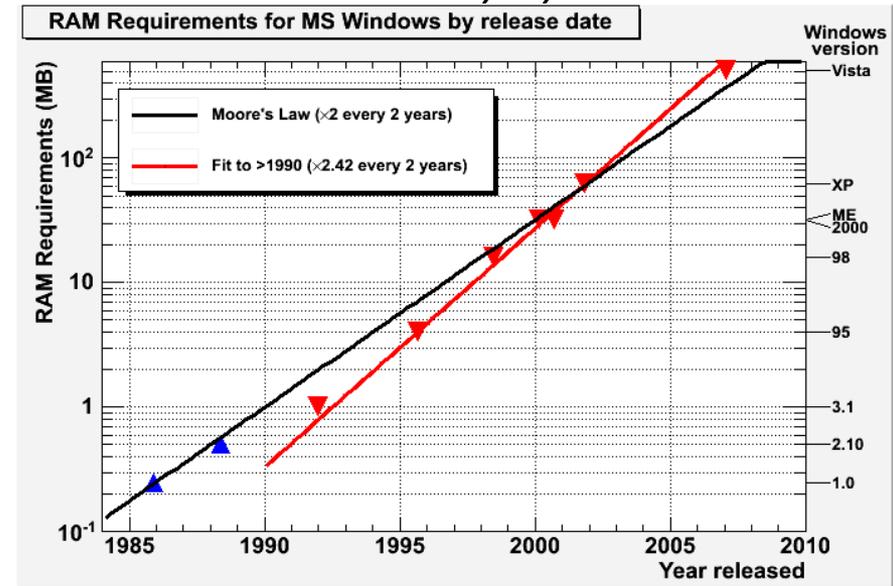
trend: $\div 2.5$ every 2 years



©2009 John C. McCallum
thanks to gnuplot

Memory Requirements

trend: $\times 2.4$ every 2 years



David Lawrence 2010

Memory cost for a desktop computer has been roughly constant over the last 2 decades.

The cost of a desktop system has been roughly constant so this keeps memory at a constant fraction of the total cost.

Multiple cores + memory

Multi-core processors are already here and commonly used. Industry has signaled that this will be the trend for the next several years. Consequence: Parallelism is required

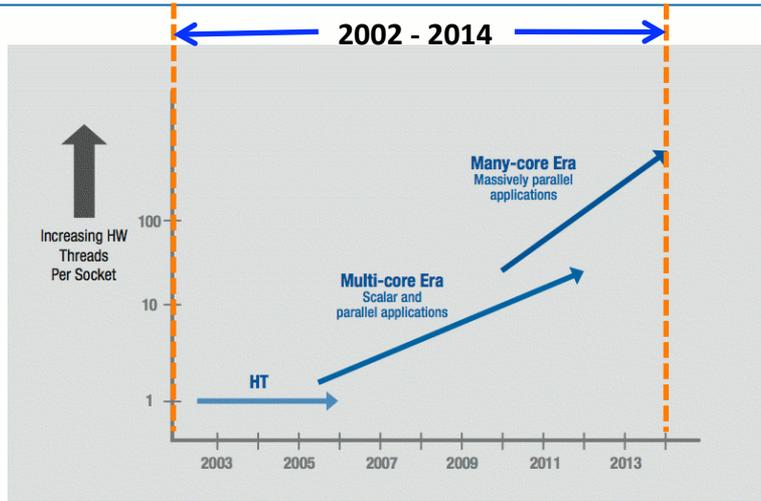
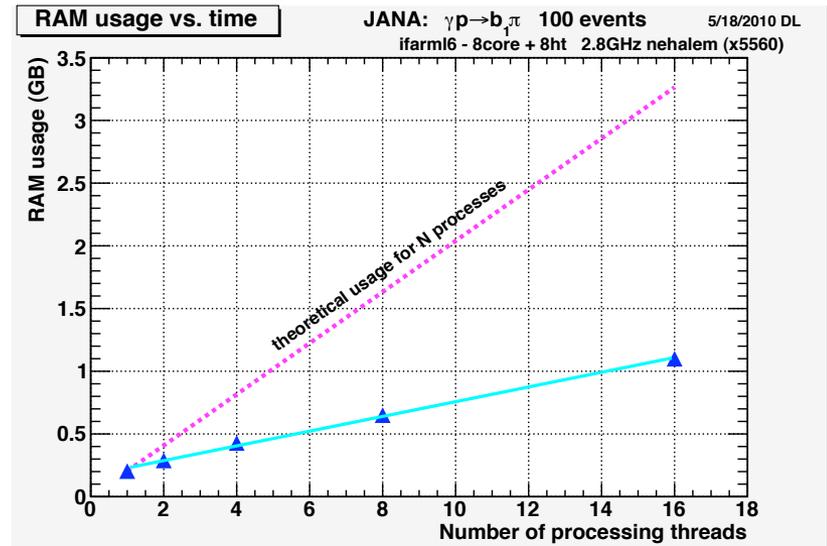
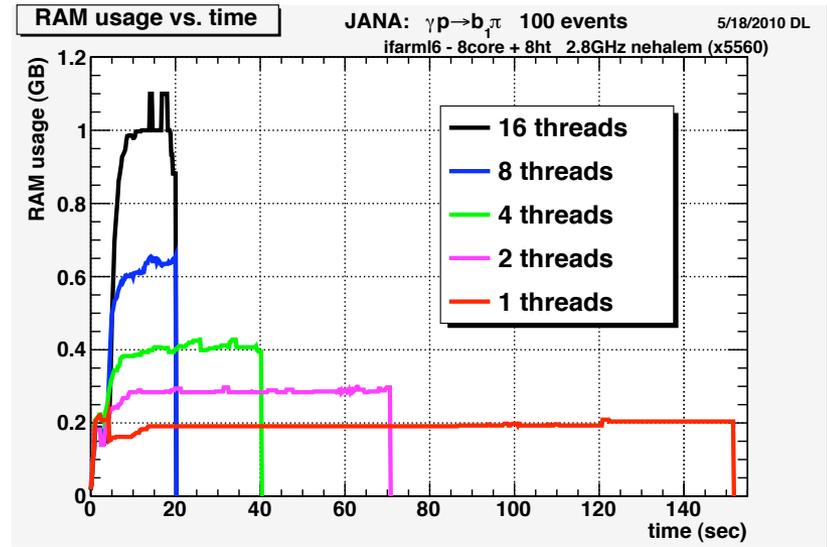


Figure 1: Current and expected eras of Intel® processor architectures

Maintaining a fixed memory capacity per core will become increasingly expensive due to limitations on the number of controllers that can be placed on a single die (#pins).

Prediction is that number of cores in the “Many-core Era” will increase faster than Moore’s law adding to the difficulty in maintaining a fixed memory capacity per core.



top

Memory usage is between 3 GB and 4 GB for single process running with 48 processing threads

CPU completely utilized in user space (like we want it!)

Negligible time spent sleeping in mutex locks or system calls

```
top - 10:31:21 up 6 days, 20:16, 2 users, load average: 24.20, 8.58, 3.13
Tasks: 662 total, 2 running, 660 sleeping, 0 stopped, 0 zombie
Cpu(s): 99.9%us, 0.0%sy, 0.0%ni, 0.0%id, 0.0%wa, 0.0%hi, 0.0%si, 0.0%st
Mem: 65980312k total, 10034708k used, 55945604k free, 161628k buffers
Swap: 1052248k total, 0k used, 1052248k free, 6610748k cached
```

PID	USER	PR	NI	VIRT	RES	SHR	S	%CPU	%MEM	TIME+	COMMAND
6298	davidl	15	0	4216m	2.8g	21m	R	4797.9	4.5	15:48.97	hd_ana
6191	davidl	15	0	19452	1624	880	R	0.7	0.0	0:00.39	top
1	root	15	0	10344	684	568	S	0.0	0.0	0:05.70	init
2	root	RT	-5	0	0	0	S	0.0	0.0	0:00.21	migration/0
3	root	34	19	0	0	0	S	0.0	0.0	0:00.16	ksoftirqd/0
4	root	RT	-5	0	0	0	S	0.0	0.0	0:00.00	watchdog/0
5	root	RT	-5	0	0	0	S	0.0	0.0	0:00.06	migration/1
6	root	34	19	0	0	0	S	0.0	0.0	0:00.00	ksoftirqd/1
7	root	RT	-5	0	0	0	S	0.0	0.0	0:00.00	watchdog/1
8	root	RT	-5	0	0	0	S	0.0	0.0	0:00.05	migration/2
9	root	34	19	0	0	0	S	0.0	0.0	0:00.00	ksoftirqd/2
10	root	RT	-5	0	0	0	S	0.0	0.0	0:00.00	watchdog/2
11	root	RT	-5	0	0	0	S	0.0	0.0	0:00.04	migration/3
12	root	34	19	0	0	0	S	0.0	0.0	0:00.00	ksoftirqd/3
13	root	RT	-5	0	0	0	S	0.0	0.0	0:00.00	watchdog/3
14	root	RT	-5	0	0	0	S	0.0	0.0	0:00.63	migration/4
15	root	34	19	0	0	0	S	0.0	0.0	0:00.29	ksoftirqd/4

I/O Scaling

Multiple processes simultaneously reading and writing to the same local disk will cause the disk head to thrash, ultimately leading to an I/O bottleneck

Multiple threads will stream events from a single file leading to much less competition for the head position

